

EC'24 Tutorial on Transaction Fee Mechanism Design

Organizers:

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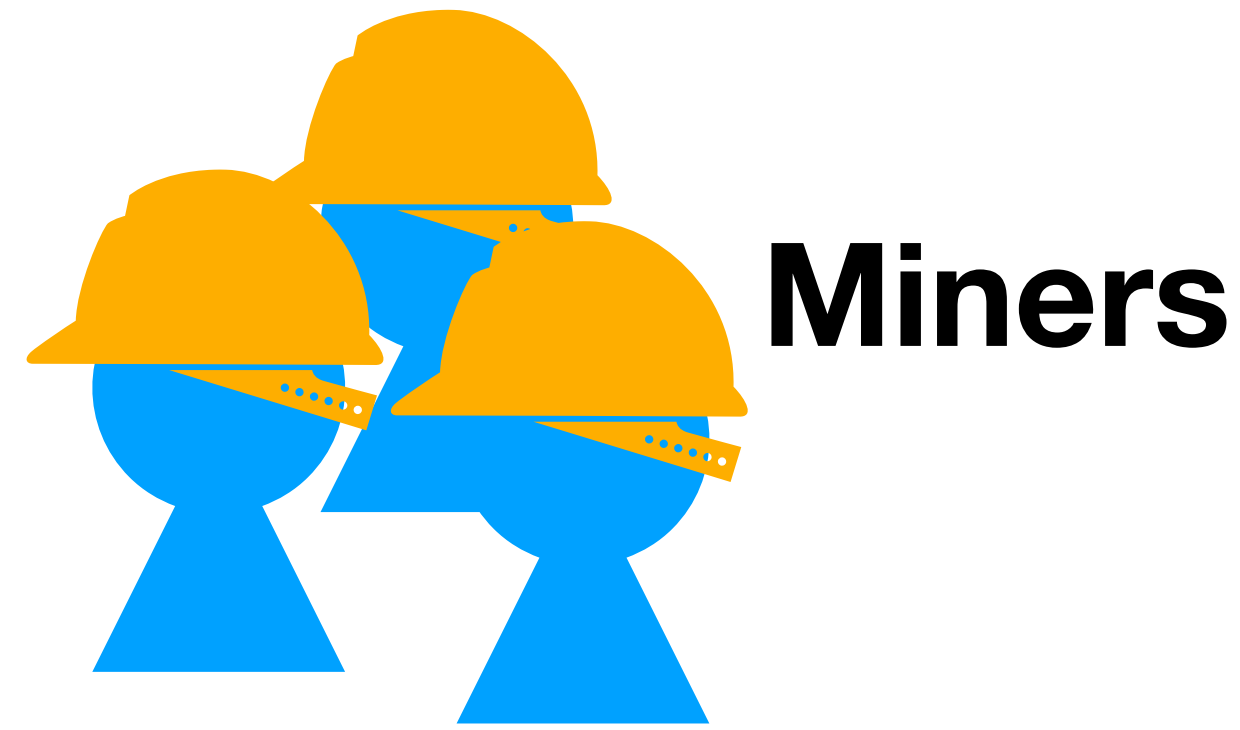
Agenda

- Lecture 1 (20 mins): TFMs for a single block
- Lecture 2 (20 mins): Dynamics TFMs
 - Break (30 mins)
- Lecture 3 (20 mins): Extensions to the TFM frameworks
- Panel discussion (30 mins):
 - Mallesh M. Pai (Rice University and Consensys)
 - Tim Roughgarden (Columbia University and a16z crypto)
 - Noam Nisan (Hebrew University of Jerusalem and Starkware)

Lecture 1: TFMs for a Single Block

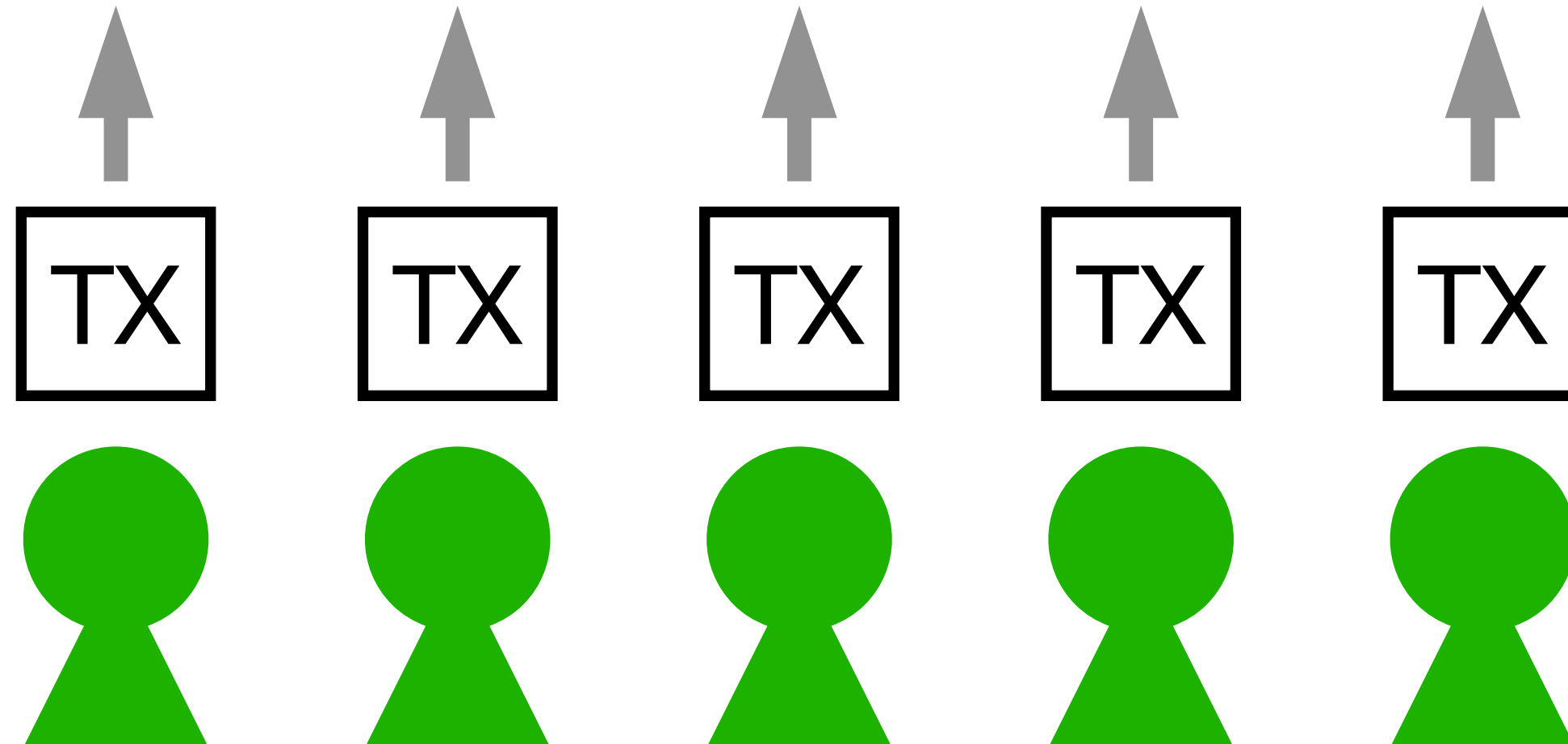
- What are blockchains and TFMs?
- TFM's desiderata
- Limitations in the single-block setting
- What can cryptography do for TFM design?

Blockchain is a public computer





A blockchain is a public computer maintained by multiple **miners**




Users



Users operate the computer by sending **transactions**

A transaction can be as simple as coin transfer

 Block:  20120974 1 Block Confirmation

 Timestamp:  17 secs ago (Jun-18-2024 08:17:47 PM +UTC) |  Confirmed within 30 secs

 Transaction Action:  Transfer 4.19705488 ETH To `0x0f967c884545d1b295aEf0281eE49688CA7255a4`

 Sponsored:

 From: `0xFd90a4bF5892dA15F863e8C385A789e583F2117D` 

 To: `0x0f967c884545d1b295aEf0281eE49688CA7255a4` 

— A transaction can also be a complex program

```
// SPDX-License-Identifier: GPL-3.0
pragma solidity >=0.7.0 <0.9.0;
/// @title Voting with delegation.
contract Ballot {
    // This declares a new complex type which will
    // be used for variables later.
    // It will represent a single voter.
    struct Voter {
        uint weight; // weight is accumulated by delegation
        bool voted;  // if true, that person already voted
        address delegate; // person delegated to
        uint vote;    // index of the voted proposal
    }

    // This is a type for a single proposal.
    struct Proposal {
        bytes32 name;    // short name (up to 32 bytes)
        uint voteCount; // number of accumulated votes
    }
```


— A transaction can also be a complex program

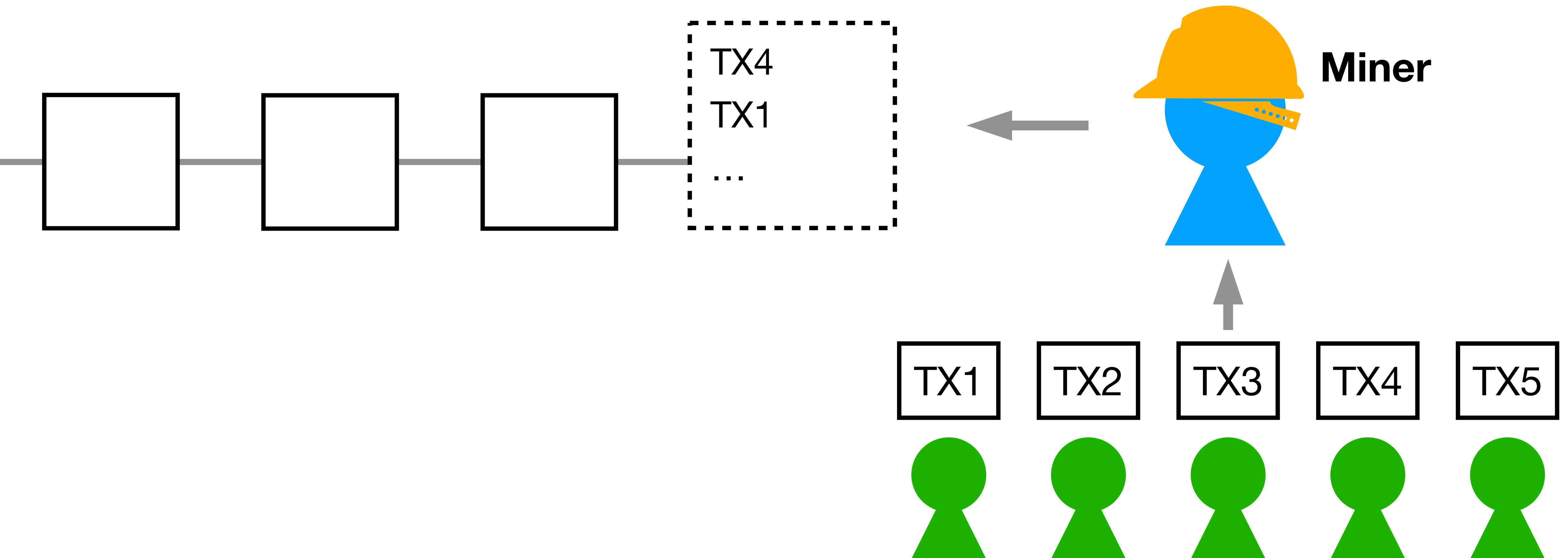
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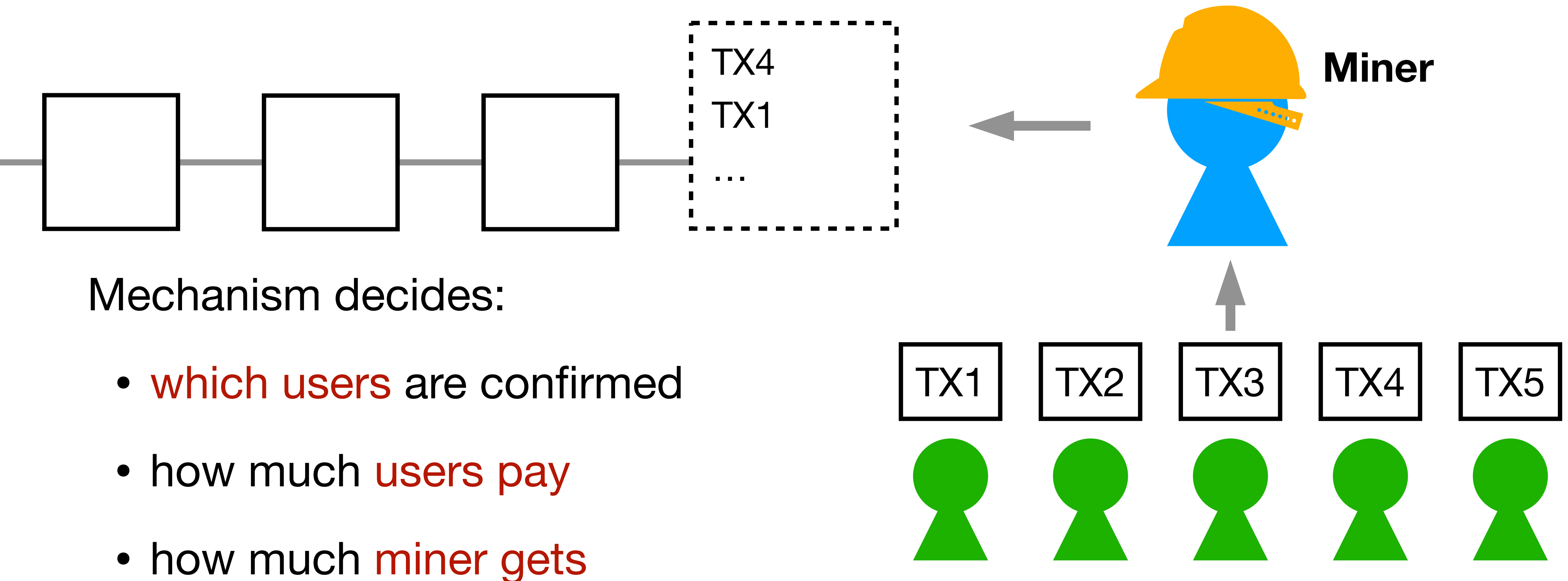
programs on blockchains are
known as “smart contracts”

— This computer is updated block-by-block

The miner packs a bunch of transactions into a **block**



Transaction fee mechanism is like an auction



What counts as good TFM's?

Bitcoin: first price auction



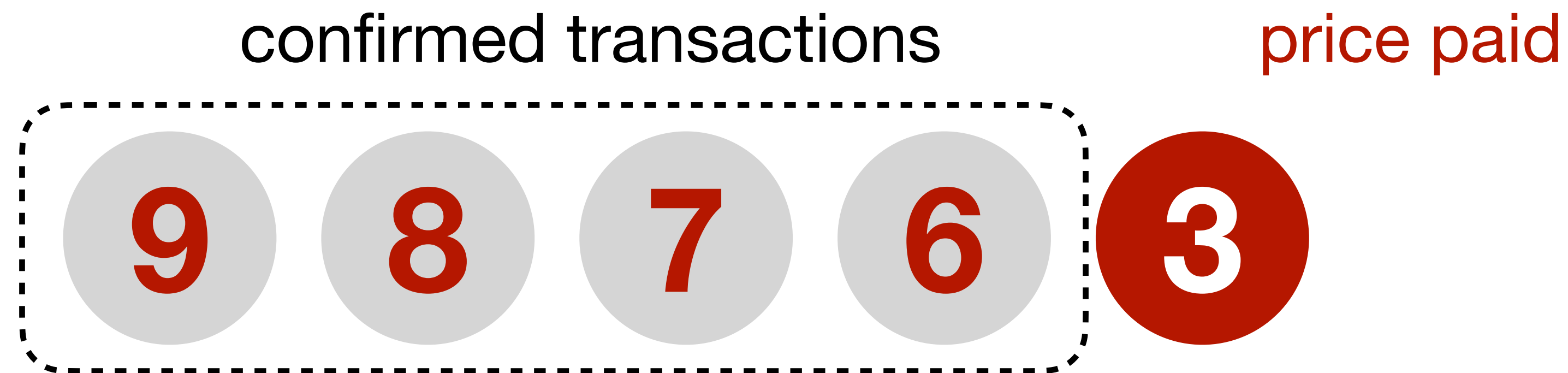
1. Top k bids are confirmed
2. Pay your own bid
3. All payment goes to the miner



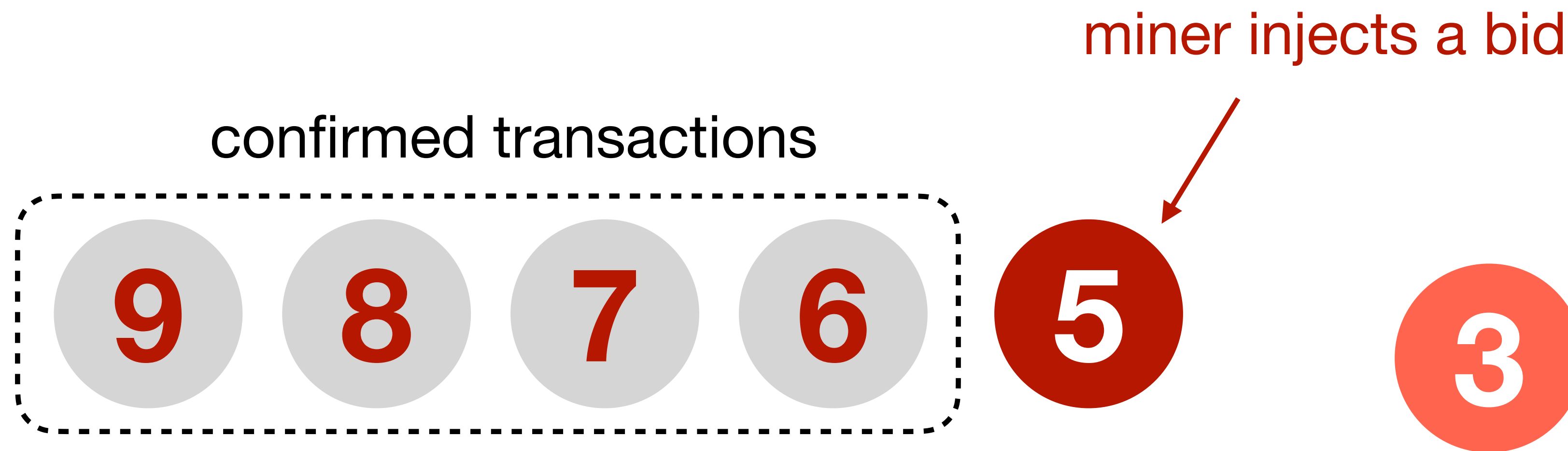
Encourage **untruthful bidding**

Truthful bidding by classical mechanisms

We can run 2nd price auction



Classical Mechanisms Fail!



Three desired properties

UIC (user incentive compatibility)

- A user's best strategy is to bid truthfully

MIC (miner incentive compatibility)

- Miner's best strategy is to implement the mechanism honestly

c-SCP (c-side-contract-proofness)

- A coalition of the miner and at most c users doesn't want to deviate

**New challenges
in decentralized context!**



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**New challenges
in decentralized context!**



!

2nd price auction is UIC, but not MIC and 1-SCP

1st price auction is MIC and c-SCP, but not UIC

Ethereum's EIP-1559 achieves all properties assuming infinite block size

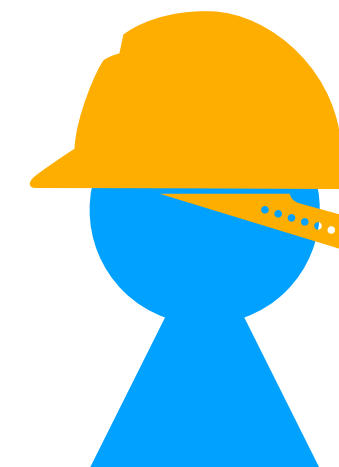
Uncongested \implies posted-price auction

- All bids \geq posted price r are confirmed, and pay r
- miner gets nothing; all payment is burnt

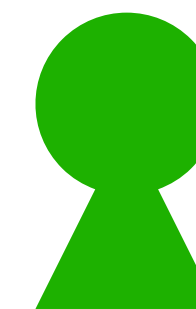


Congested \implies first price auction

Without burning, miner-user coalition can bypass the price r



bid 5,
cash back 3



$r = 5$,
but value = 3

— Dream mechanism is impossible!

Theorem

Suppose the block size is finite.

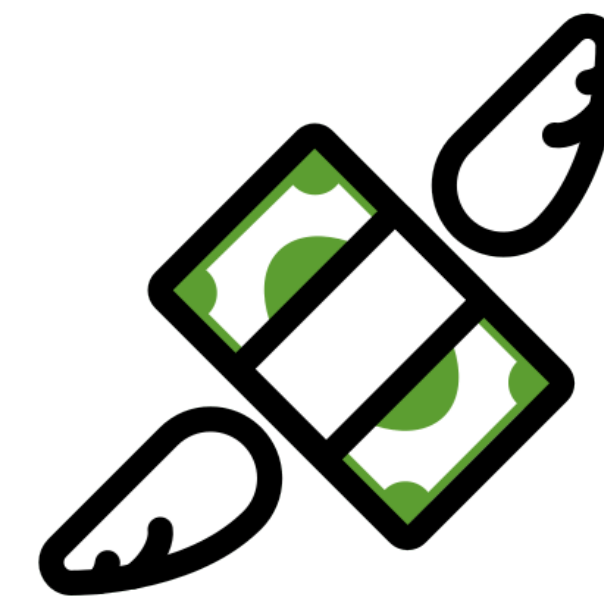
No non-trivial TFM can satisfy UIC and 1-SCP at the same time.

Zero miner revenue is inherent

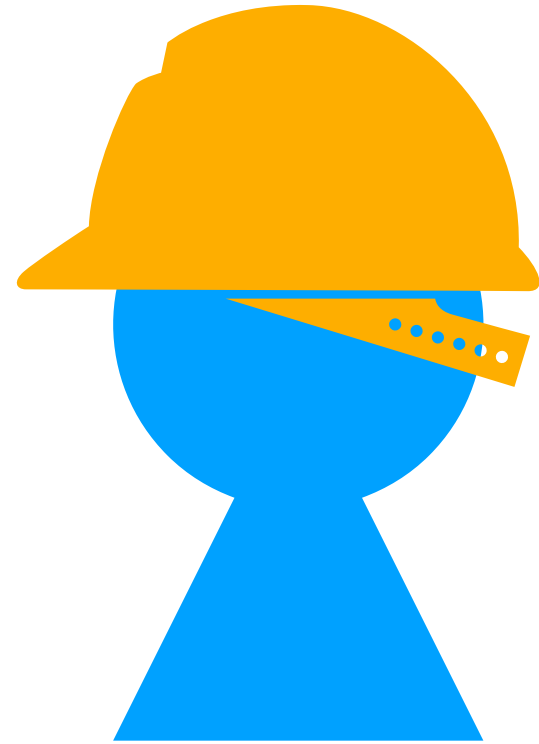
Theorem

For any TFM that satisfies UIC and 1-SCP,
miner revenue must be zero.

Burning in EIP1559 is necessary!

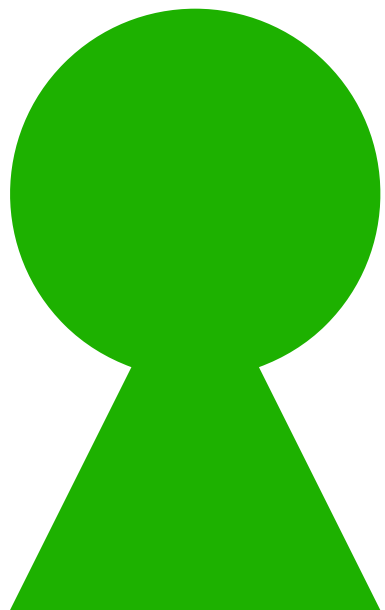


Strategy Space in Plain Model



After seeing others' bids, a **miner** can

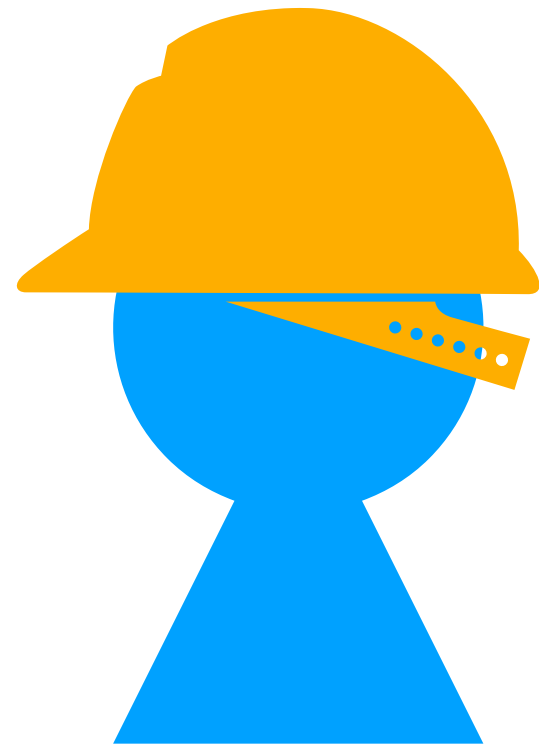
- inject fake bids
- create a block arbitrarily



After seeing others' bids, a **user** can

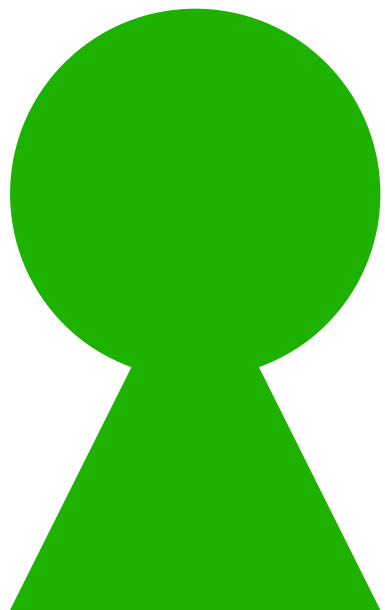
- bid untruthfully
- inject fake bids

Strategy Space in MPC-assisted Model



~~After seeing others' bids,~~ a **miner** can

- inject fake bids
- ~~create a block arbitrarily~~



~~After seeing others' bids,~~ a **user** can

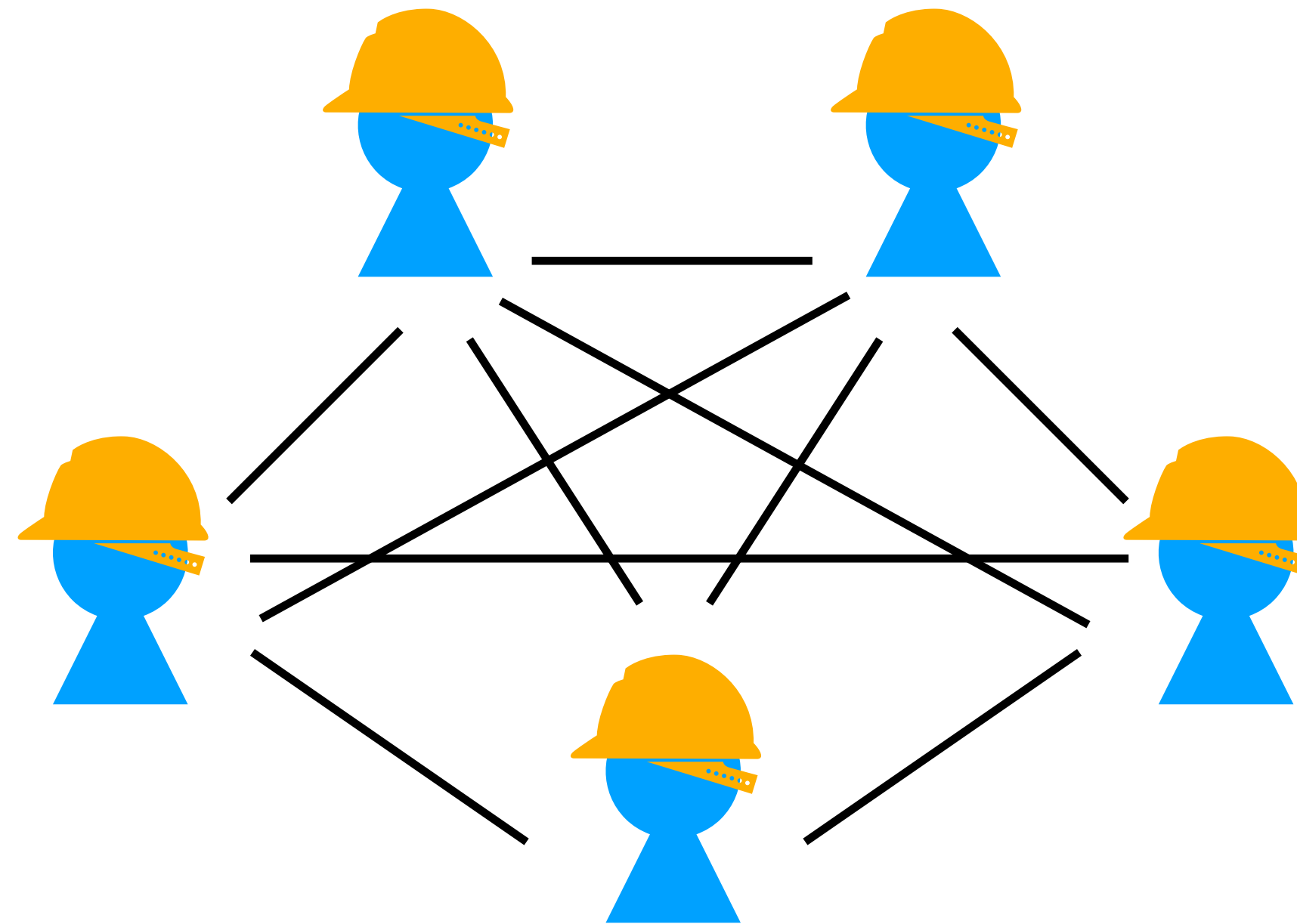
- bid untruthfully
- inject fake bids

Posted-price with random selection

- All bids \geq **posted-price** r are **eligible**
- Randomly choose k eligible bids to confirm
- Each confirmed bid pays r
- All payments are **burnt**

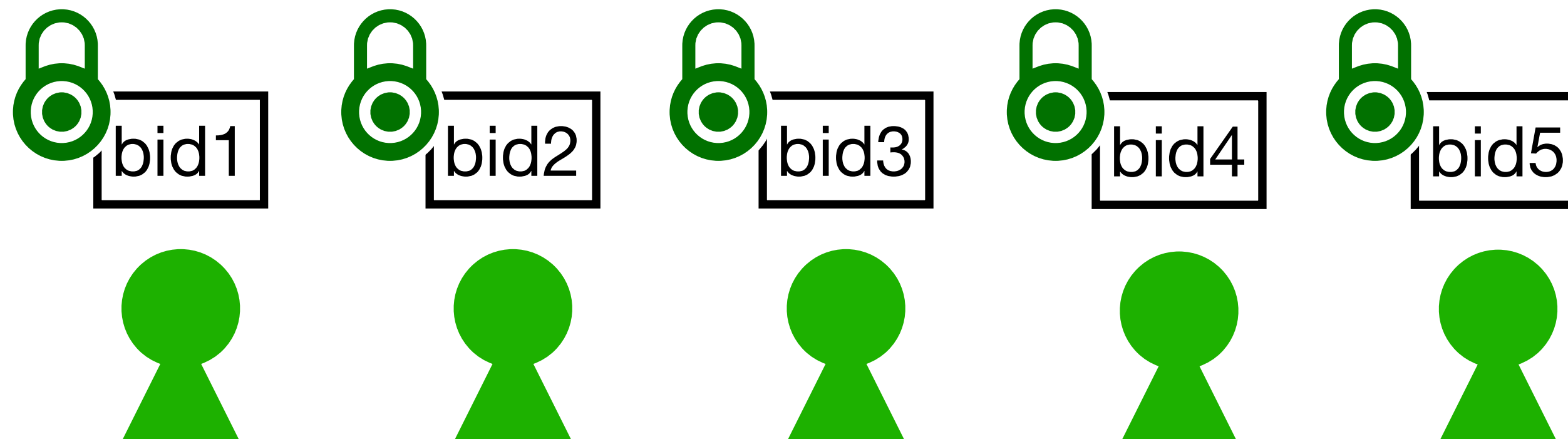
This mechanism is **UIC + MIC + 1-SCP**

MPC-Assisted Model



Multiple miners jointly run **multiparty computation (MPC)**

A user **secret-shares** its bid,
and sends each share to each miner



Take away from lecture 1

- A blockchain is a public computer
- Transaction fee mechanisms (TFMs) allocate block space
- New design feature: burning
- New challenge: miner and miner-user deviation
- In plain model: UIC + 1-SCP \implies trivial mechanism

Some simplifications in lecture 1

- Focus on a **single** block
 - In practice: multiple block in the long term (lecture 2)
- All transactions have **equal size**
 - In practice: different size (“gas” model in Ethereum)
- Transaction **order** does not matter in the block
 - In practice: order matters! (lecture 3)
- A single miner fully controls one block
 - Depend on protocols, e.g. MPC-assisted mechanism (end of this lecture) or proposer-builder separation (lecture 3)

Thank you!



— backup: Ethereum's EIP-1559

base fee tip

- Each bid specifies (r, t)
- All bids \geq **base-fee** r are **eligible**
- Miner confirms up to k eligible bids with highest tips
- Each confirmed bid pays $r + t$
- Miner gets all the tips, the base fee is **burnt**

