

EC'24 Tutorial on Transaction Fee Mechanism Design

Matheus V. X. Ferreira



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Agenda

- Lecture 1 (20 mins): TFMs for a single block
- **Lecture 2 (20 mins): Dynamics TFMs**
 - Break (30 mins)
- Lecture 3 (20 mins): Extensions to the TFM frameworks
- Panel discussion (30 mins):
 - Malleh M. Pai (Rice University and Consensys)
 - Tim Roughgarden (Columbia University and a16z crypto)
 - Noam Nisan (Hebrew University of Jerusalem and Starkware)

Why dynamic mechanism?

- **User experience:** users might be willing to wait for future block space
- **[Example]** In each round, a second-price auction allocates a single item for sale.
 - Alice is patient and her value is 6.
 - The first block of the bids is {4, 5}.
 - What should Alice bid?
 - Second block the bids are {1, 2}.
 - If she would be equally happy with waiting for the second block she overpaid by 3.

Related work (more on the tutorial website)

- Dynamical Analysis of the EIP-1559 Ethereum Fee Market, 2021 [S. Leonardos, B. Monnot, D. Reijsbergen, E. Skoulakis, G. Piliouras]
- Dynamic Posted-Price Mechanisms for the Blockchain Transaction-Fee Market, 2021 [M.V.X. Ferreira, D.J. Moroz, D.C. Parkes, M. Stern]
- Dynamic Transaction Fee Mechanism Design, 2024 [M. Pai, M. Resnick]
- Serial Monopoly on Blockchains, 2023 [N. Nisan]
- Competitive Revenue Extraction from Time-Discounted Transactions in the Semi-Myopic Regime, 2024 [Y. Gafni, A. Yaish]

Digital Transaction

- Intermediaries: Visa, Mastercard, Amex, ACH
- Great user experience, but arguably expensive for merchants



\$5



\$0.10 fee

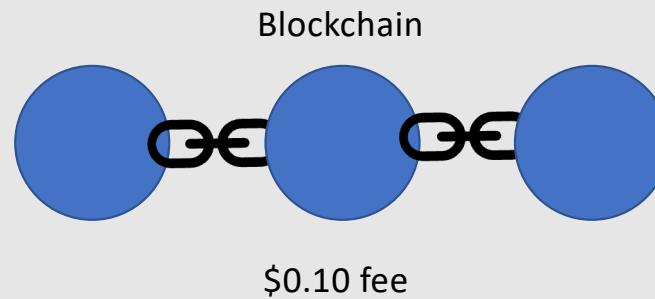
Transaction outcome:

- Rejected (insufficient funds, etc...) or confirmed at \$0.10

Digital Transaction on Blockchains



\$5



Transaction outcome:

- Rejected (insufficient funds, etc...), confirmed at \$0.10, or delayed:
 - Either eventually confirmed for \$0.10, or
 - It cannot abort and you increase the fee to \$0.15 to speed up the confirmation.

Model

- In each round
 - N “identical” items for sale
 - A new set of buyers arrive
 - A mechanism allocates the items to at most N buyers
 - Any unallocated buyer stays for the next round (or **timeout**)

Challenge: The mechanism designer has weak commitment power because multiple “pseudonymous” miners/sequencers/proposers take turns to implement the allocation rule.

Objective

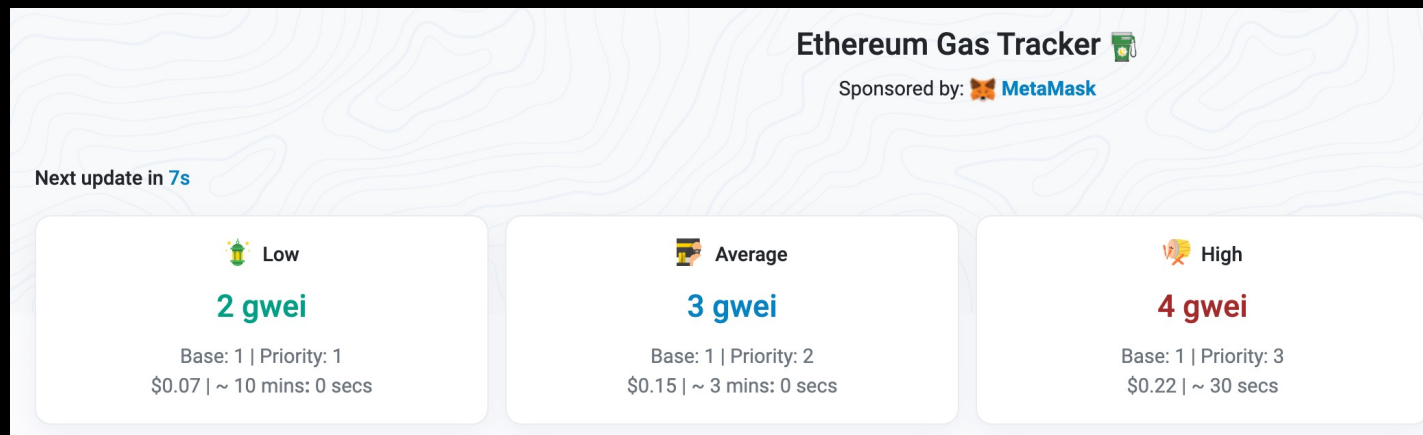
Maximize Welfare
(i.e., allocate space
to who benefits the
most)

Simplicity
(e.g., Incentive
Compatible)

Posted-price provides simplicity



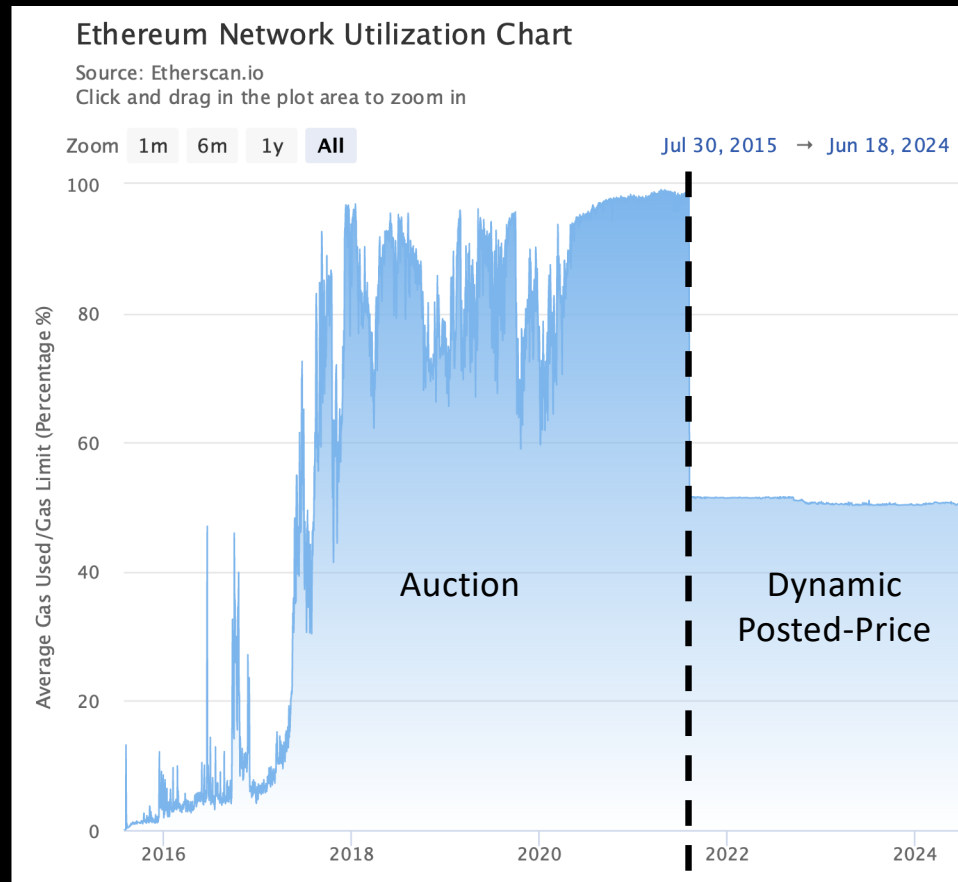
Simplicity ...



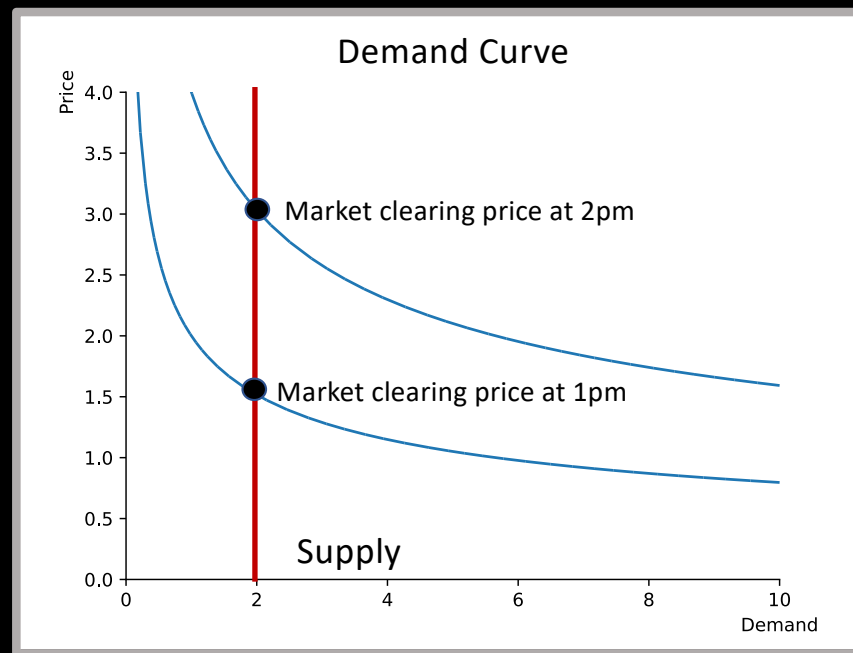
Source: [3 Gwei](#) | [Ethereum Gas Tracker](#) | [Etherscan](#)

Accessed: 06/19/2024

Welfare



How to dynamically price block space given future demand is unknown?



Approach (EIP 1559 [Buterin et al., '19])

1. Each block contains a posted-price: $Price_t$.
2. Miner can **ONLY** include transactions with a bid above $Price_t$.
3. Bidder pays $Price_t$.
4. Compute the posted-price for next block using a pricing rule

Pending Transactions				
TX ID	Sender	Receiver	Value	Budget
01	Alice	Bob	10	1
02	Charlie	David	15	0.01
03	Bob	Charlie	1	0.5

Pricing Rules

Utilization-based
(EIP-1559)

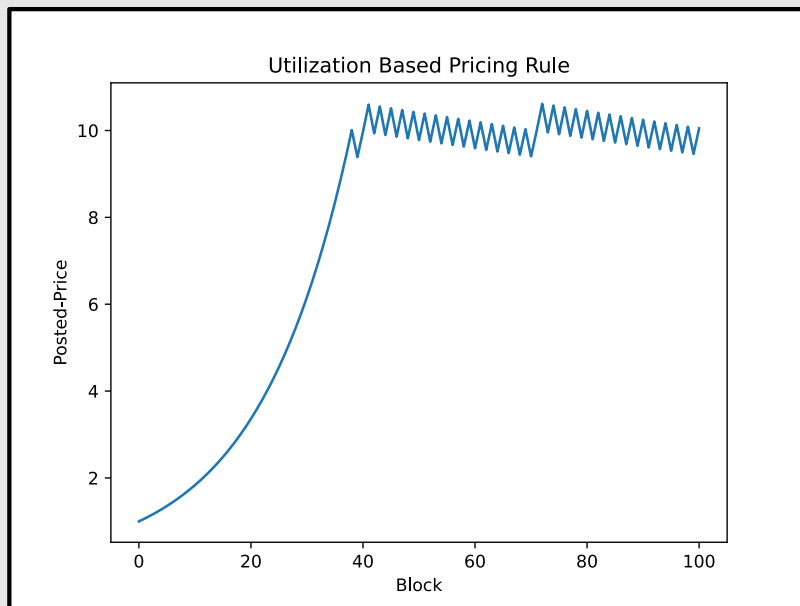
Utilization-based rule

$$Price_{t+1} = Price_t (1 + \alpha(Utilization - Target))$$

- Block ***Utilization*** = $\frac{\# Transactions\ in\ Block}{Block\ Capacity}$.
- **Target** utilization ($Target = 1/2$ in EIP-1559).

Instability of utilization-based rule [FMPS '21]

- Suppose 50 slots for sale and 100 users that bid 10.



$$\begin{aligned} Price_{t+1} &= Price_t(1 + \alpha(Utilization - Target)) \\ &= 1/2 \\ &= Price_t \left(1 \mp \frac{\alpha}{2}\right) \end{aligned}$$

Welfare-based pricing rule [FMPS '21]

$$Welfare(Block) = \sum_{i \in Block} v_i$$

$$Price_{t+1} = \alpha \frac{Welfare(Block)}{Capacity} + (1 - \alpha) Price_t$$

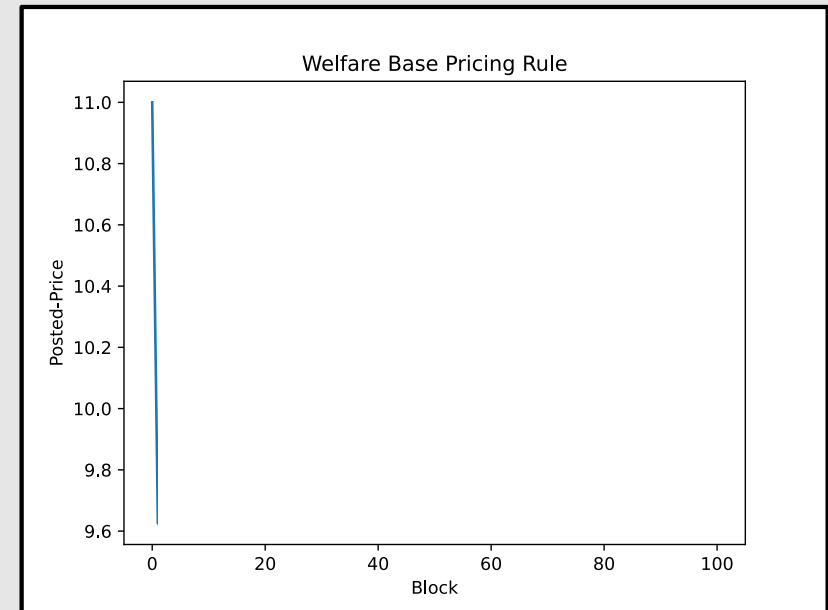
- Each transaction contributes $\frac{\alpha v_i}{Capacity}$ (where v_i is the bid of bidder i).

Example

- Consider 50 slots for sale and 100 users that bid 10.
- Case 1 ($Price_t^W > 10$):

$$\begin{aligned} Price_{t+1}^W &= \alpha \frac{Welfare(Block)}{Capacity} + (1 - \alpha) Price_t^W \\ &= (1 - \alpha) Price_t^W \\ &< Price_t^W \end{aligned}$$

- Then eventually $Price_{t+1}^W \leq 10$.

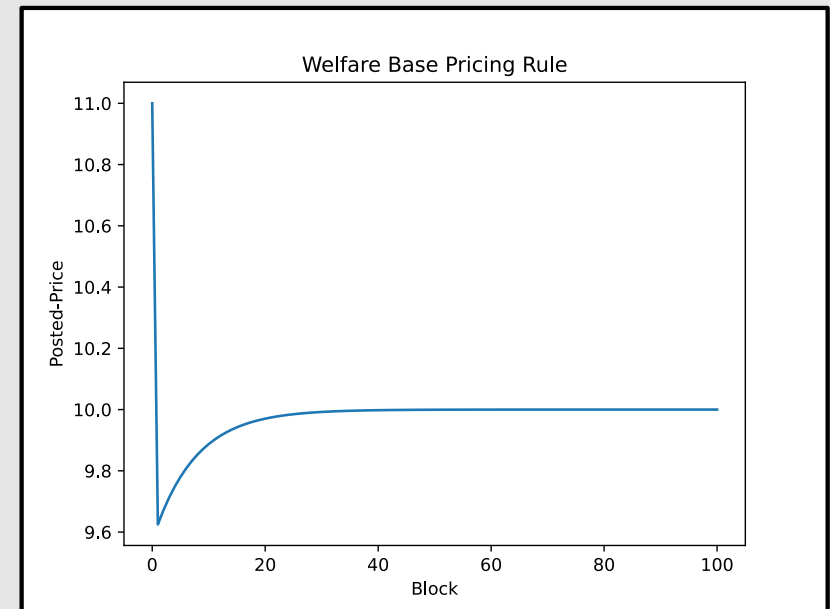


Example: Welfare-based

- Case 2 ($Price_t^W \leq 10$):

$$\begin{aligned} Price_{t+1}^W &= \alpha \frac{Welfare(Block)}{Capacity} + (1 - \alpha) Price_t^W \\ &= 10\alpha + (1 - \alpha) Price_t^W \\ &\geq Price_t^W \text{ and } \leq 10 \end{aligned}$$

- Thus, sequence of prices is monotone increasing.
- From monotone convergence, posted-price converge.



Why burn the posted-price?

- Without burning
 - If $value < price$, seller asks buyer to bid $price$
 - If buyer wins, buyer pays $price$ and seller refunds $price - value$
 - Considered undesirable, but it improves welfare
- With burning
 - Seller does not benefit from including buyers with $value < price$
 - Seller does benefit from deviating when $value > price$ (e.g, impose a reserve price if they have Bayesian beliefs about values)
 - Example: over 90% of miners deviate from EIP-1559 by selling their block in the MEV-boost.

Opportunities for future work

- What is simplicity for multi-shot mechanisms? (e.g., no regret?)
- AI-assisted tooling (e.g., Ethereum gas tracker)
- In this talk slots are identical. In practice, slots are not identical

Value = (1st slot value, 2nd slot value, ...)

- Result in the posted-price being always smaller than 1st slot clearing price (congested)
- Is Welfare maximization a good objective? What about fairness (MEV)?
[Ferreira, Parkes '23]

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Thank you